



# MayBMS: A System for Managing Large Uncertain and Probabilistic Databases

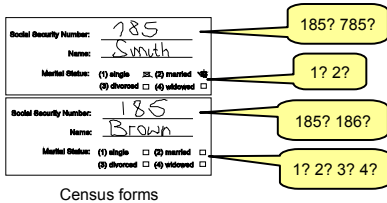
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## 1. Motivation

Goal: manage uncertain information in different application scenarios: data integration, scientific data collections, census ...



There are  $2 * 2 * 2 * 4 = 32$  possible instances of the forms information!

## Features of MayBMS

- ✓ scalable DBMS for supporting uncertain and probabilistic data
- ✓ purely relational representation of attribute-level uncertainty
- ✓ efficient query processing
- ✓ query language for probabilistic databases

## 2. U-relational Databases

$R[SSN, N, MS]$ : personal information

$U_{R[SSN]}$	TID	SSN	$V_1 \rightarrow D_1$	$U_{R[MS]}$	TID	MS	$V_1 \rightarrow D_1$	$U_{R[N]}$	TID	N	$V_1 \rightarrow D_1$
	$t_1$	185	$x \rightarrow 1$		$t_1$	1	$z \rightarrow 1$		$t_1$	Smith	
	$t_1$	785	$x \rightarrow 2$		$t_1$	2	$z \rightarrow 2$		$t_2$	Brown	
	$t_2$	185	$y \rightarrow 1$		$t_2$	1	$w \rightarrow 1$				
	$t_2$	186	$y \rightarrow 2$		$t_2$	2	$w \rightarrow 2$				
					$t_2$	3	$w \rightarrow 3$				
					$t_2$	4	$w \rightarrow 4$				

world table: prob. distribution of the variables

W	$V \rightarrow D$	Pr
	$x \rightarrow 1$	0.4
	$x \rightarrow 2$	0.6
	$y \rightarrow 1$	0.7
	$y \rightarrow 2$	0.3
	$z \rightarrow 1$	0.8
	$z \rightarrow 2$	0.2
	$w \rightarrow 1$	0.25
	$w \rightarrow 2$	0.25
	$w \rightarrow 3$	0.25
	$w \rightarrow 4$	0.25

$S[SSN, ST]$ : credit status

$U_{S[SSN, ST]}$	TID	SSN	ST	$V_1 \rightarrow D_1$
	$s_1$	185	bad	$w \rightarrow 3$
	$s_1$	185	good	$w \rightarrow 4$

Construct a possible world: pick a value for each variable

R	SSN	Name	MS
$t_1$	785	Smith	2
$t_2$	186	Brown	4

S	SSN	ST
$s_1$	185	good

Probability of the world:  $0.6 * 0.3 * 0.2 * 0.25 = 0.009$

encode attribute alternatives and correlations with variables

## 3. Query Language

### World-set Algebra

✓ extend relational algebra with uncertainty-specific constructs e.g.:

- **conf**: confidence computation
- **repair by key**: create the possible repairs of an instance violating a key constraint
- **assert**: remove worlds violating a constraint
- ✓ semantics: evaluate the query in each world
- ✓ properties
  - generic: independent from representation details
  - conservative over relational algebra: right degree of expressive power
  - efficient evaluation: simple encoding of positive relational algebra + possible into positive relational algebra queries on U-relational databases

## 4. Query Evaluation

a) possible( $\pi_N(\sigma_{MS=3}(R))$ )

Query on U-relational databases:

$$\pi_N(\sigma_{MS=3}(U_{R[SSN]} \bowtie_{\varphi \& \psi} U_{R[MS]}))$$

merge( $\pi_{N,R}, \pi_{M,S}$ )

$\varphi = I.TID = r.TID$

$\psi = (I.V_1 = r.V_1 \rightarrow I.D_1 = r.D_1) \& (I.V_2 = r.V_2 \rightarrow I.D_2 = r.D_2)$

b) repair-key $_{SSN}(R)$

$U_{R[SSN]}$	TID	SSN	$V_1 \rightarrow D_1$	$V_2 \rightarrow D_2$
	$t_1$	185	$x \rightarrow 1$	$x_1 \rightarrow 1$
	$t_1$	785	$x \rightarrow 2$	
	$t_2$	185	$y \rightarrow 1$	$x_1 \rightarrow 2$
	$t_2$	186	$y \rightarrow 2$	

new variable for each non-unique SSN value

← ensure consistent variable assignment

c)  $\pi_{MS}(R \bowtie_{SSN} (\sigma_{ST=bad}(S)))$

Query on column-stores

$$\pi_{MS}(\text{merge}(\pi_{SSN,R}, \pi_{MS,S}) \bowtie_{SSN} \sigma_{ST=bad}(S)) \equiv \pi_{MS}(\text{merge}(\pi_{\emptyset}(\pi_{SSN,R} \bowtie_{SSN} \sigma_{ST=bad}(S)), \pi_{MS,S}))$$

push merge up

intermediate result:

$U_{TII}$	TID	$V_1 \rightarrow D_1$	$V_2 \rightarrow D_2$
	$t_1, s_1$	$x \rightarrow 1$	$w \rightarrow 3$
	$t_2, s_1$	$y \rightarrow 1$	$w \rightarrow 3$

final result:

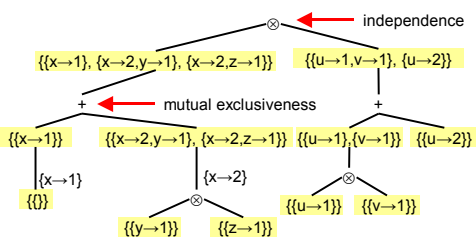
$U_{P[MS]}$	TID	MS	$V_1 \rightarrow D_1$	$V_2 \rightarrow D_2$	$V_3 \rightarrow D_3$
	$t_1, s_1$	1	$x \rightarrow 1$	$w \rightarrow 3$	$z \rightarrow 1$
	$t_1, s_1$	2	$x \rightarrow 2$	$w \rightarrow 3$	$z \rightarrow 2$
	$t_2, s_1$	3	$y \rightarrow 1$	$w \rightarrow 3$	

## 5. Confidence Computation

$U_{R[A]}$	A	$V_1 \rightarrow D_1$	$V_2 \rightarrow D_2$
	1	$x \rightarrow 1$	
	1	$x \rightarrow 2$	$y \rightarrow 1$
	1	$x \rightarrow 2$	$z \rightarrow 1$
	1	$u \rightarrow 1$	$v \rightarrow 1$
	1	$u \rightarrow 2$	

confidence of (A:1) = probability of the world-set defined by

$\{\{x \rightarrow 1\}, \{x \rightarrow 2, y \rightarrow 1\}, \{x \rightarrow 2, z \rightarrow 1\}, \{u \rightarrow 1, v \rightarrow 1\}, \{u \rightarrow 2\}\}$



## 6. Experiments

- ✓ extended TPC-H population generator 2.6 to generate U-relational databases
- ✓ parameters: scale (s), uncertainty ratio (x), correlation ratio (z), max alternatives per field (m), drop after correlation (p)
- ✓ each generated world has the sizes of relations and join selectivities of the original TPC-H one-world case
- ✓ queries translated into SQL and run on PostgreSQL

s	z	TPC-H dbsize	#worlds	dbsize	#worlds	dbsize
0.5	0.1	853	$10^{13588}$	3843	$10^{17496}$	5427
0.5	0.5	853	$10^{25288}$	3856	$10^{27340}$	6682
1	0.1	1706	$10^{12610}$	7683	$10^{18496}$	11264
1	0.5	1706	$10^{12610}$	7712	$10^{18496}$	13312
		X=0.0	X=0.001		X=0.1	

Fig. 1: Number of worlds and size in MB of the U-relational db for different scale, uncertainty and correlation ratios.

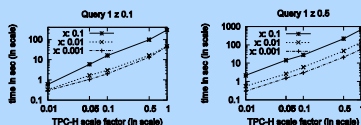


Fig. 2: Query evaluation

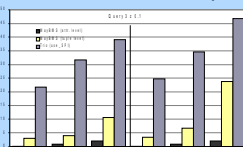


Fig. 3: Attribute- vs. tuple-level representation

## 7. Selected Publications

1. Fast and simple relational processing of uncertain data, ICDE'08
2. Conditioning Probabilistic Databases, Technical report '08
3. MayBMS: Managing incomplete information with probabilistic world-set decompositions, ICDE'07, Demo paper
4. 10<sup>10</sup> worlds and beyond: efficient representation and processing of incomplete information, ICDE'07
5. From complete to incomplete information and back, SIGMOD'07